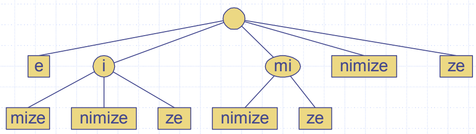
Tries



Outline and Reading

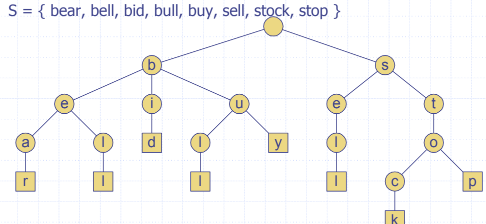
- Standard tries (§9.2.1)
- Compressed tries (§9.2.2)
- Suffix tries (§9.2.3)
- Huffman encoding tries (§9.3.1)

Preprocessing Strings

- Preprocessing the pattern speeds up pattern matching queries
  - After preprocessing the pattern, KMP's algorithm performs pattern matching in time proportional to the text size
- If the text is large, immutable and searched for often (e.g., works by Shakespeare), we may want to preprocess the text instead of the pattern
- A trie is a compact data structure for representing a set of strings, such as all the words in a text
  - A tries supports pattern matching queries in time proportional to the pattern size

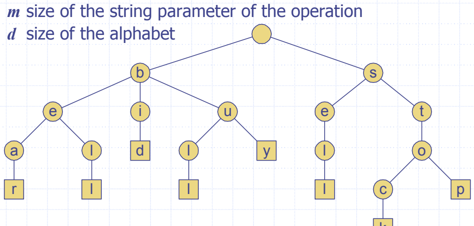
Standard Trie (1)

- The standard trie for a set of strings S is an ordered tree such that:
  - Each node but the root is labeled with a character
  - The children of a node are alphabetically ordered
  - The paths from the external nodes to the root yield the strings of S
- Example: standard trie for the set of strings
  S = { bear, bell, bid, bull, buy, sell, stock, stop }



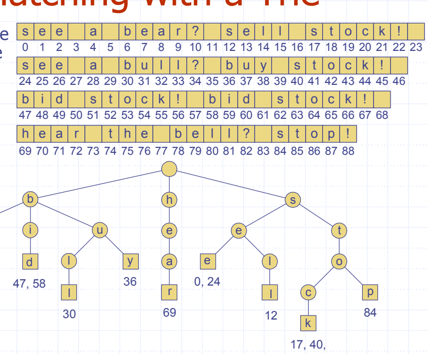
Standard Trie (2)

- A standard trie uses $O(n)$ space and supports searches, insertions and deletions in time $O(dm)$, where:
  - n total size of the strings in S
  - m size of the string parameter of the operation
  - d size of the alphabet



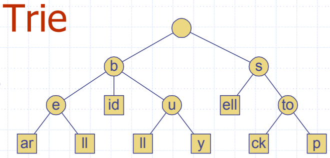
Word Matching with a Trie

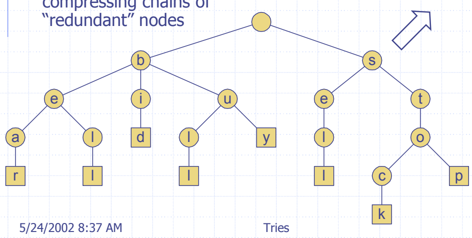
- We insert the words of the text into a trie
- Each leaf stores the occurrences of the associated word in the text



Compressed Trie

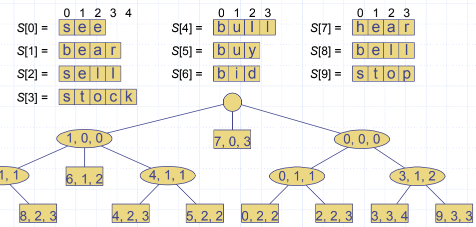
- A compressed trie has internal nodes of degree at least two
- It is obtained from standard trie by compressing chains of "redundant" nodes





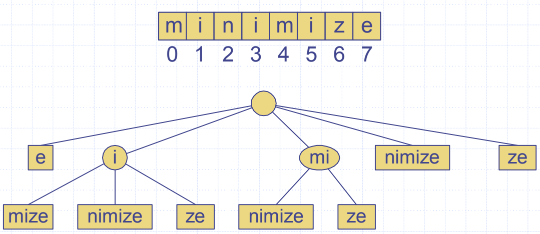
Compact Representation

- Compact representation of a compressed trie for an array of strings:
  - Stores at the nodes ranges of indices instead of substrings
  - Uses $O(s)$ space, where s is the number of strings in the array
  - Serves as an auxiliary index structure



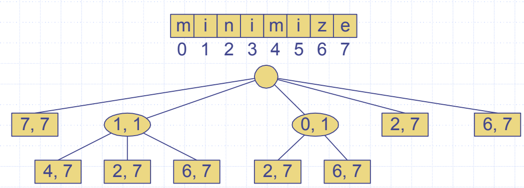
Suffix Trie (1)

- The suffix trie of a string X is the compressed trie of all the suffixes of X



Suffix Trie (2)

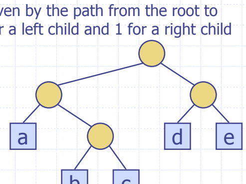
- Compact representation of the suffix trie for a string X of size n from an alphabet of size d
  - Uses $O(n)$ space
  - Supports arbitrary pattern matching queries in X in $O(dm)$ time, where m is the size of the pattern



Encoding Trie (1)

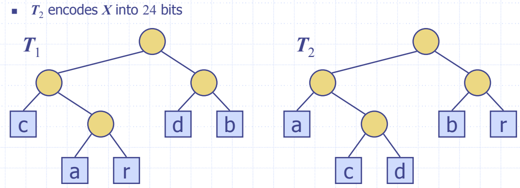
- A code is a mapping of each character of an alphabet to a binary code-word
- A prefix code is a binary code such that no code-word is the prefix of another code-word
- An encoding trie represents a prefix code
  - Each leaf stores a character
  - The code word of a character is given by the path from the root to the leaf storing the character (0 for a left child and 1 for a right child

00	010	011	10	11
a	b	c	d	e



Encoding Trie (2)

- Given a text string X, we want to find a prefix code for the characters of X that yields a small encoding for X
  - Frequent characters should have long code-words
  - Rare characters should have short code-words
- Example
  - X = abracadabra
  - T_1 encodes X into 29 bits
  - T_2 encodes X into 24 bits



Huffman's Algorithm

- Given a string X, Huffman's algorithm construct a prefix code the minimizes the size of the encoding of X
- It runs in time $O(n + d \log d)$, where n is the size of X and d is the number of distinct characters of X
- A heap-based priority queue is used as an auxiliary structure

```
Algorithm HuffmanEncoding(X)
  Input string X of size n
  Output optimal encoding trie for X
  C ← distinctCharacters(X)
  computeFrequencies(C, X)
  Q ← new empty heap
  for all c ∈ C
    T ← new single-node tree storing c
    Q.insert(getFrequency(c), T)
  while Q.size() > 1
    f1 ← Q.minKey()
    T1 ← Q.removeMin()
    f2 ← Q.minKey()
    T2 ← Q.removeMin()
    T ← join(T1, T2)
    Q.insert(f1 + f2, T)
  return Q.removeMin()
```


Example

X = abracadabra

Frequencies

a	b	c	d	r
5	2	1	1	2